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## EFFICIENT BLOCK PRECONDITIONING FOR A $\mathbb{C}^1$ FINITE ELEMENT DISCRETISATION OF THE DIRICHLET BIHARMONIC PROBLEM

J. PESTANA†\$, R. MUDDLE†‡, M. HEIL†, F. TISSEUR†\$, AND M. MIHAJLOVIƇ

Abstract. We present an efficient block preconditioner for the two-dimensional biharmonic Dirichlet problem discretised by  $C^1$  bicubic Hermite finite elements. In this formulation each node in the mesh has four different degrees of freedom (DOFs). Grouping DOFs of the same type together leads to a natural blocking of the Galerkin coefficient matrix. Based on this block structure, we develop two preconditioners: a  $2 \times 2$  block diagonal preconditioner (BD) and a block bordered diagonal (BBD) preconditioner. We prove mesh-independent bounds for the spectra of the BD-preconditioned Galerkin matrix under certain conditions. The eigenvalue analysis is based on the fact that the proposed preconditioner, like the coefficient matrix itself, is symmetric positive definite and is assembled from element matrices. We demonstrate the effectiveness of the inexact version of the BBD preconditioner, which exhibits near-optimal scaling in terms of computational cost with respect to the discrete problem size. Finally, we study robustness of this preconditioner with respect to domain distortion and element stretching.

**Key words.** biharmonic equation; Hermite bicubic finite elements; block preconditioning; conjugate gradient method; algebraic multigrid

AMS subject classifications. 65F08, 65F10, 65N22

1. Introduction. The biharmonic operator is a key component in mathematical models of a number of important physical problems. It arises in plane strain and plane stress elasticity problems, where the solution is expressed in terms of an Airy stress function [25, p. 79], [30, p. 288], and in plate bending problems. It also occurs in the stream-function-vorticity formulation of Stokes flow [22].

The strong formulation of the biharmonic problem seeks the function  $u \in C^4(\Omega)$  that satisfies

$$\nabla^4 u = f \tag{1.1}$$

in the domain  $(x_1, x_2) \in \Omega \subset \mathbb{R}^2$  with piecewise smooth boundary  $\partial\Omega$  and source function  $f \in L_2(\Omega)$  subject to the Dirichlet boundary conditions

$$u = g_1, \quad \frac{\partial u}{\partial \hat{n}} = g_2 \text{ on } \partial \Omega,$$
 (1.2)

where  $\frac{\partial u}{\partial \hat{n}}$  denotes the outward normal derivative and  $g_1$  and  $g_2$  are given functions. In the context of the plate bending problem the case  $g_1 = g_2 = 0$  corresponds to a clamped boundary.

Numerical schemes for solving (1.1)–(1.2) either approach the problem directly, or reformulate it as a mixed formulation (i.e., solve a system of two second-order problems). The advantages of using the former approach include better asymptotic accuracy for the same level of grid resolution [1, Theorem 5.4], [7, Theorems 6.1.6 and

<sup>&</sup>lt;sup>†</sup>School of Mathematics, The University of Manchester, Manchester, M13 9PL, UK, (jennifer.pestana@manchester.ac.uk, rlmuddle@googlemail.com, m.heil@maths.man.ac.uk, francoise.tisseur@manchester.ac.uk)

<sup>&</sup>lt;sup>‡</sup>School of Computer Science, The University of Manchester, Manchester, M13 9PL, UK (milan@cs.man.ac.uk)

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7.1.6 and p. 392], and a symmetric positive definite coefficient matrix for the discrete problem. Conversely, for the mixed formulation, discretisation (by a finite difference or finite element method, for example) results in a linear algebraic system that is symmetric, but indefinite. (Efficient preconditioners for such systems are based on block partitioning of the coefficient matrix and the application of multigrid to the Dirichlet Laplacian blocks [24], or to the Schur complement system [14].)

In this paper we consider a conforming  $C^1$  finite element approach [21], for which the standard weak form is to find  $u \in H^2(\Omega)$  satisfying (1.2) such that

$$\int_{\Omega} \nabla^2 u \, \nabla^2 v \, d\Omega = \int_{\Omega} f v \, d\Omega, \tag{1.3}$$

holds for all test functions  $v \in H_0^2(\Omega)$ , where  $H_0^2(\Omega) = \{v \in H^2(\Omega) \mid v = \frac{\partial v}{\partial \hat{n}} = 0 \text{ on } \partial \Omega\}$ . The discrete weak formulation is obtained by restricting (1.3) to a finite dimensional space  $S(\Omega) \subset H^2(\Omega)$ , for which we adopt a basis associated with the bicubic Hermite (Bogner-Fox-Schmit) finite elements [3, p. 72]; these are formed from a tensor product of one-dimensional Hermite polynomials. The  $C^1$  continuity across element boundaries is ensured by assigning four degrees of freedom (DOF) to each node, corresponding to four different basis functions [3, p. 72].

The finite element approximation of (1.3) is then obtained by solving a linear system  $\mathcal{A}x = \mathbf{b}$ , where  $\mathcal{A} \in \mathbb{R}^{N \times N}$  is a large, sparse and symmetric positive definite (SPD) matrix and  $\mathbf{b} \in \mathbb{R}^N$ . Such systems are usually solved by iterative methods, with the conjugate gradient (CG) method a popular choice [13, Ch. 2]. Grouping together the unknowns corresponding to the same DOF type leads to the following natural  $4 \times 4$  blocking of the coefficient matrix

$$\mathcal{A} = \begin{bmatrix} A_{11} & A_{12} & A_{13} & A_{14} \\ A_{12}^T & A_{22} & A_{23} & A_{24} \\ A_{13}^T & A_{23}^T & A_{33} & A_{34} \\ A_{14}^T & A_{24}^T & A_{34}^T & A_{44} \end{bmatrix}, \tag{1.4}$$

where  $A_{ij} \in \mathbb{R}^{n \times n}$ ,  $i, j = 1, \ldots, 4$ , and N = 4n, where n is the number of interior nodes. Since the biharmonic operator is fourth order, the two-norm condition number of the matrix  $\mathcal{A}$  behaves as  $\kappa(\mathcal{A}) = O(h^{-4})$ , where h is the mesh parameter (assuming uniform discretisation). This rapid growth in the condition number has a detrimental effect on the convergence speed of the CG method. This problem can be rectified by effective preconditioning.

There are a number of preconditioning strategies for a conforming  $C^1$  discretisations of (1.1)–(1.2). The proposed methods include additive Schwartz methods [12], [33], [34], BPX preconditioning [21], Steklov-Poincare operator-based preconditioning [19], problem-specific multigrid methods [4], [6], [15], [26], [31], and fast auxiliary space (FASP) preconditioning [32].

In this paper we propose two novel preconditioners that are fully algebraic and are assembled from the element matrices in an analogous manner to the matrix  $\mathcal{A}$ , making them both simple to implement. The first of these preconditioners is a  $2 \times 2$  block diagonal (BD) matrix. The positive definiteness of  $\mathcal{A}$  and the assembly of the preconditioner from element matrices means that analysis based on the general ideas of Wathen [27], [28], [29] can be applied to demonstrate that mesh-independent convergence is guaranteed in certain cases.

The second preconditioner introduced in this paper is a computationally cheaper block bordered diagonal (BBD) approximation of the block diagonal preconditioner, that is feasible for larger problems, and that can be implemented in a cost-effective manner. For this second preconditioner we provide some spectral analysis. We then employ numerical experiments to demonstrate mesh independent convergence rates and show that it is possible to deploy off-the-shelf multigrid approximations for certain matrix blocks. Robust convergence is confirmed by a number of case studies.

The paper is organised as follows. In Section 2 we discuss the finite element assembly process of the matrix  $\mathcal{A}$  in (1.4) and relevant aspects of the conjugate gradient method. Section 3 describes the new block diagonal preconditioner. We characterise the eigenvalues of the preconditioned matrix and give conditions for mesh-independent convergence. However, the preconditioner is costly to apply within the preconditioned conjugate gradient method. We therefore introduce a more practical block bordered diagonal preconditioner in Section 4, and provide an eigenvalue analysis. We propose an inexact version of the BBD preconditioner, which involves matrix lumping and algebraic multigrid approximation. The computational cost of applying this preconditioner is shown to scale linearly with the problem size N, thus giving an optimal Krylov solver for the biharmonic system. Finally, we present numerical experiments in Section 5 that verify the effectiveness of the inexact BBD preconditioner and investigate its robustness with respect to changes in the domain and element shape.

- 2. Preliminaries. In this section we describe the details of the finite element assembly process for the biharmonic problem and introduce the preconditioned conjugate gradient (PCG) method.
- **2.1. The finite element assembly process.** The analysis of the spectra of the preconditioned matrices in later sections will be based on the fact that the finite element matrix  $\mathcal{A}$  in (1.4) is assembled from element contributions. In this section we describe this assembly process.

We discretise (1.3) using  $C^1$  Hermite finite elements, defined in a reference domain with local co-ordinates  $(s_1, s_2) \in \overline{\Omega} = [-1, 1]^2$ . The solution within the element is represented as

$$u = \sum_{j=1}^{4} \sum_{k=1}^{4} U_{jk} \,\bar{\psi}_{jk}(s_1, s_2),$$

where  $U_{jk}$  are the unknown coefficients and  $\bar{\psi}_{jk}$  are the reference Hermitian basis functions. The subscript j represents the node number and k enumerates the DOF type, such that at node j,  $U_{jk}$  interpolates u,  $\frac{\partial u}{\partial s_1}$ ,  $\frac{\partial u}{\partial s_2}$  and  $\frac{\partial^2 u}{\partial s_1 \partial s_2}$  for  $k = 1, \ldots, 4$ , respectively. The same basis functions are used to isoparametrically map the reference element to the actual element  $\Omega_e$ .

Consider now a finite element discretisation of the domain  $\Omega$  consisting of M elements and let  $A_e \in \mathbb{R}^{16 \times 16}$ , e = 1, ..., M, be the biharmonic element matrices associated with these elements. The matrices  $A_e$  are symmetric positive semi-definite and each entry is of the form

$$(A_e)_{ij} = \int_{\overline{\Omega}} \nabla^2 \bar{\psi}_{i_1 i_2} \nabla^2 \bar{\psi}_{j_1 j_2} |J_e| \, d\Omega,$$

where  $i=4(i_2-1)+i_1$  and  $j=4(j_2-1)+j_1$  and  $J_e$  is the element Jacobian. Consequently, multiplying  $A_e$  by a vector  $\boldsymbol{u}\in\mathbb{R}^{16}$ , with elements  $u_j=u_{j_1j_2}$ , is

equivalent to computing integrals of linear combinations of basis vectors, that is,

$$(A_e \boldsymbol{u})_i = \int_{\overline{\Omega}} \nabla^2 \bar{\psi}_{i_1 i_2} \left( \sum_{j_1, j_2 = 1}^4 \nabla^2 \left( u_{j_1 j_2} \bar{\psi}_{j_1 j_2} \right) \right) d\Omega.$$

Thus, the nullspace vectors of  $A_e$  can be thought of in terms of linear combinations of certain basis functions. These nullspace basis functions are harmonic functions, i.e., functions for which the Laplacian is zero (see (1.3)). It is straightforward to verify that a basis for these harmonic functions is

1, 
$$s_1$$
,  $s_2$ ,  $s_1s_2$ ,  $s_1^2 - s_2^2$ ,  $s_2(s_1^2 - s_2^2/3)$ ,  $s_2(s_1^2/3 - s_2^2)$  and  $s_1s_2(s_1^2 - s_2^2)$ , (2.1)

from which the nullspace of  $A_e$  can be computed.

Now let us describe the assembly process of (1.4) mathematically. We introduce the matrix  $L_e \in \mathbb{R}^{16 \times N}$  that maps the entries of  $A_e$  to entries of A. Then

$$\mathcal{A} = \sum_{i=1}^{M} L_e^T A_e L_e = L^T \operatorname{diag}(A_e) L \in \mathbb{R}^{N \times N}, \tag{2.2}$$

where

$$L = \begin{bmatrix} L_1^T & L_2^T & \dots & L_M^T \end{bmatrix}^T \in \mathbb{R}^{N \times N}$$
 (2.3)

and  $\operatorname{diag}(A_e)$  is a block diagonal matrix with each block consisting of an element matrix  $A_e$ , i = 1, ..., M. The matrix  $\operatorname{diag}(A_e)$  is related to the differential operator and the choice of basis functions, while L provides information about the geometry and boundary conditions.

During this assembly process, unknowns corresponding to the same DOF type are grouped together and this leads to the natural blocking of the coefficient matrix as:

$$\mathcal{A} = \begin{bmatrix} A_{11} & A_{12} & A_{13} & A_{14} \\ A_{12}^T & A_{22} & A_{23} & A_{24} \\ A_{13}^T & A_{23}^T & A_{33} & A_{34} \\ A_{14}^T & A_{24}^T & A_{34}^T & A_{44} \end{bmatrix} \cdot \begin{array}{c} u \\ \frac{\partial u}{\partial s_1} \\ \frac{\partial u}{\partial s_2} \\ \frac{\partial^2 u}{\partial s_1 \partial s_2} \\ \frac{\partial^2 u}{\partial s_2 \partial s_2} \\ \end{array}$$
(2.4)

The unknown vector  $\boldsymbol{x}$  and the right-hand side  $\boldsymbol{b}$  are blocked accordingly.

**2.2.** The conjugate gradient method. The conjugate gradient method (CG) is perhaps the best known Krylov subspace method for solving sparse linear systems, and is suitable for systems with an SPD coefficient matrix. The relative error  $e^{(k)} = x - x^{(k)}$  after k iterations of CG is bounded by [13, p. 51]

$$\frac{\|\boldsymbol{e}^{(k)}\|_{\mathcal{A}}}{\|\boldsymbol{e}^{(0)}\|_{\mathcal{A}}} \le 2\left(\frac{\sqrt{\kappa(\mathcal{A})} - 1}{\sqrt{\kappa(\mathcal{A})} + 1}\right)^{k},$$

where  $\kappa(\mathcal{A}) = \lambda_{\text{max}}(\mathcal{A})/\lambda_{\text{min}}(\mathcal{A})$  is the two-norm condition number of  $\mathcal{A}$ . As mentioned in the introduction and verified numerically in Table 2.1, the condition number of the matrix  $\mathcal{A}$  in (1.4) is  $\kappa(\mathcal{A}) = O(h^{-4})$  and this suggests a significant deterioration in convergence speed of the CG solver as the mesh is refined (see the computations in Section 5).

Table 2.1

Extremal eigenvalues and 2-norm condition number of  $\mathcal{A}$  from (1.4) as a function of the problem size N.

Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$
$\overline{N}$	36	196	900	3844	15876
$\lambda_{\min}$	56.20	18.45	4.94	1.26	0.32
$\lambda_{ m max}$	1287	5705	23399	94179	377295
$\kappa$	23	309	4735	74912	$1.20 \times 10^{6}$

The problem of ill-conditioning can be alleviated by solving an equivalent preconditioned system  $\mathcal{P}^{-\frac{1}{2}}\mathcal{A}\mathcal{P}^{-\frac{1}{2}}\boldsymbol{y} = \mathcal{P}^{-\frac{1}{2}}\boldsymbol{b}$  with  $\boldsymbol{x} = \mathcal{P}^{-\frac{1}{2}}\boldsymbol{y}$ , where  $\mathcal{P} \in \mathbb{R}^{N \times N}$  is SPD. Note that the CG algorithm itself requires only a linear system solve with  $\mathcal{P}$ at each iteration, i.e., the matrix  $\mathcal{P}^{-\frac{1}{2}}$  is never explicitly formed. The error of the preconditioned CG iterates can be bounded by

$$\frac{\|\boldsymbol{e}^{(k)}\|_{\mathcal{A}}}{\|\boldsymbol{e}^{(0)}\|_{\mathcal{A}}} \leq 2\left(\frac{\sqrt{\kappa(\mathcal{P}^{-1}\mathcal{A})}-1}{\sqrt{\kappa(\mathcal{P}^{-1}\mathcal{A})}+1}\right)^{k}.$$

The error bound shows that the convergence of the CG method is accelerated when the condition number of the preconditioned matrix  $\mathcal{P}^{-1}\mathcal{A}$  is small. It can also be shown that fast convergence rates are achieved when the eigenvalues belong to a small number of tightly bounded clusters (see, for example, [13, Section 3.1]). If the eigenvalues of  $\mathcal{P}^{-1}\mathcal{A}$  can be bounded independently of the mesh size h (and possibly other problem parameters) then  $\mathcal{P}$  is an optimal preconditioner, in the sense of convergence of the conjugate gradient method. If, in addition, linear systems involving  $\mathcal{P}$  can e solved in a manner that scales linearly with the problem size then we have an optimal solver.

**3. An ideal preconditioner.** We first consider the block diagonal preconditioner

$$\mathcal{P}_{BD} = \begin{bmatrix} A_{11} & A_{12} & A_{13} \\ A_{12}^T & A_{22} & A_{23} \\ A_{13}^T & A_{23}^T & A_{33} \\ & & & A_{44} \end{bmatrix} . \tag{3.1}$$

Since any principal submatrix of an SPD matrix  $\mathcal{A}$  is iteself SPD [17, p. 397], the preconditioner  $\mathcal{P}_{BD} \in \mathbb{R}^{N \times N}$  is also symmetric positive definite. Additionally,  $\mathcal{P}_{BD}$  is formed from a subset of the block matrices  $A_{ij}$  of  $\mathcal{A}$  and so it is possible to assemble  $\mathcal{P}_{BD}$  from the element matrix contributions in a manner analogous to that described in Section 2.1. Thus,

$$\mathcal{P}_{BD} = L^T \operatorname{diag}(P_e)L, \tag{3.2}$$

where  $P_e$  is obtained from  $A_e$ , with values that would be assembled into  $A_{i4}$  or  $A_{i4}^T$  set to zero for i=1,2,3. The element contribution to the preconditioner (henceforth, the element preconditioner)  $P_e$ , like  $A_e$ , is symmetric positive semidefinite, but it has rank 11 rather than 8. Straightforward computation shows that the nullspace of  $P_e$  is spanned by vectors corresponding to 1,  $s_1$ ,  $s_2$ ,  $s_1^2 - s_2^2$  and  $s_1^3(2s_2 - 1) + s_2^3(1 - 2s_1) + 3s_1s_2(s_2 - s_1)$ . Note that the last of these functions is a combination of the last three functions in (2.1). Consequently, the nullspace of diag( $P_e$ ) is contained in the nullspace of diag( $P_e$ ) as stated in the following lemma.

LEMMA 3.1. Let diag $(P_e)$  be as in (3.2) and diag $(A_e)$  be as in (2.2). Then  $\operatorname{null}(P_e) \subset \operatorname{null}(A_e)$  and  $\operatorname{null}(\operatorname{diag}(P_e)) \subset \operatorname{null}(\operatorname{diag}(A_e))$ .

This result will be relevant in the subsequent analysis.

We next investigate analytically the spectral properties of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$ . For convenience we introduce the notation

$$\mathcal{A} = \begin{bmatrix} A_{11} & A_{12} & A_{13} & A_{14} \\ A_{12}^T & A_{22} & A_{23} & A_{24} \\ A_{13}^T & A_{23}^T & A_{33} & A_{34} \\ \hline A_{14}^T & A_{24}^T & A_{34}^T & A_{44} \end{bmatrix} = \begin{bmatrix} A & B \\ B^T & A_{44} \end{bmatrix}, \quad \mathcal{P}_{BD} = \begin{bmatrix} A \\ A_{44} \end{bmatrix}.$$
(3.3)

Then the eigenvalues of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  are characterised by the following theorem. Theorem 3.2. Assume that  $\operatorname{rank}(B) = r$  in (3.3). Then  $\mathcal{P}_{BD}^{-1}\mathcal{A}$ , with  $\mathcal{A} \in \mathbb{R}^{N \times N}$  and  $\mathcal{P}_{BD} \in \mathbb{R}^{N \times N}$  given by (1.4) and (3.1), respectively, has N-2r eigenvalues equal to 1. The remaining 2r eigenvalues  $\lambda$  satisfy

$$0 < 1 - \sqrt{\mu_{\text{max}}} \le \lambda \le 1 + \sqrt{\mu_{\text{max}}} < 2$$

where  $\mu_{\max} \in (0,1)$  is the largest eigenvalue of  $A_{44}^{-1}B^TA^{-1}B$ .

*Proof.* Since  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  is similar to  $\mathcal{P}_{BD}^{-\frac{1}{2}}\mathcal{A}\mathcal{P}_{BD}^{-\frac{1}{2}}$ , which is symmetric positive definite, any eigenvalue  $\lambda$  of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  is real and positive. Using (3.3), we see that  $\lambda$  satisfies

$$A\mathbf{u} + B\mathbf{v} = \lambda A\mathbf{u},\tag{3.4}$$

$$B^T \mathbf{u} + A_{44} \mathbf{v} = \lambda A_{44} \mathbf{v}, \tag{3.5}$$

where  $\boldsymbol{u} \in \mathbb{R}^{3n}$  and  $\boldsymbol{v} \in \mathbb{R}^n$  are not simultaneously zero and N = 4n.

If  $\lambda = 1$  then (3.4) implies that  $B\mathbf{v} = \mathbf{0}$ , i.e., that  $\mathbf{v} = \mathbf{0}$  or  $\mathbf{v} \in \text{null}(B)$ . We can find n-r linearly independent vectors in null(B) for which (3.4) and (3.5) are satisfied with u = 0. Otherwise, v = 0 and it follows from (3.5) that  $u \in \text{null}(B^T)$ . Since we can find 3n-r linearly independent vectors in null( $B^T$ ), we have that 1 is an eigenvalue of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  with multiplicity 4n-2r=N-2r.

If  $\lambda \neq 1$  then (3.4) implies that  $(\lambda - 1)^{-1}A^{-1}Bv = u$  and substituting for uin (3.5) gives that

$$A_{44}^{-1}B^TA^{-1}B\mathbf{v} = (\lambda - 1)^2\mathbf{v}.$$

From this we see that non-unit eigenvalues  $\lambda$  of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  are given by  $\lambda=1+\sqrt{\mu}$  and  $\lambda = 1 - \sqrt{\mu}$ , where  $\mu$  is a nonzero eigenvalue of  $A_{44}^{-1}B^TA^{-1}B$ . Also, since  $\mathcal{A}$  is positive definite,  $0 < \mu < 1$  [17, Theorem 7.7.7]. The result follows.  $\square$ 

The rank of B is at most n, so at least 2n eigenvalues are equal to one, while the largest non-unit eigenvalue is less than 2, regardless of the mesh size. The focus of the remainder of this section is to bound the smallest non-unit eigenvalue, since doing so ensures mesh-independent convergence.

To bound the smallest eigenvalue of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  we adapt the analysis of Wathen [27], [28], [29] to our case. The basic idea is to determine the eigenvalues of the preconditioned element matrix  $\operatorname{diag}(P_e)^{-1}\operatorname{diag}(A_e)$  and to then obtain mesh-independent bounds using Rayleigh quotients. In our case this approach is complicated by the fact that the singular matrices  $\operatorname{diag}(A_e)$  and  $\operatorname{diag}(P_e)$  have nullspaces of different dimensions. However, we can still apply the general methodology since we know from Lemma 3.1 that  $\operatorname{null}(\operatorname{diag}(P_e)) \subset \operatorname{null}(\operatorname{diag}(A_e)) \subset \mathbb{R}^{16M}$ .

To deal with the different nullspaces involved it is useful to introduce certain subspaces of  $\mathbb{R}^{16M}$ . Specifically, we define:

$$\mathcal{R} := \operatorname{range}(\operatorname{diag}(A_e)), \quad \mathcal{Z} := \operatorname{null}(\operatorname{diag}(A_e)),$$

$$\mathcal{N} := \operatorname{null}(\operatorname{diag}(P_e)), \quad \mathcal{M} := \mathcal{Z} \cap \mathcal{N}^{\perp}.$$
(3.6)

where  $\mathcal{N}^{\perp}$  is the space of all vectors orthogonal to vectors in  $\mathcal{N}$ . This decomposition satisfies  $\mathbb{R}^{16M} = \mathcal{R} + \mathcal{N} + \mathcal{M}$ , with  $\mathcal{N} \subset \mathcal{Z}$ . Additionally, since the matrices  $\operatorname{diag}(A_e)$  and  $\operatorname{diag}(P_e)$  are block diagonal, the basis vectors of  $\mathcal{R}$ ,  $\mathcal{N}$ ,  $\mathcal{Z}$  and  $\mathcal{M}$  can be constructed from their element contributions.

In addition to the spaces defined above, we require the following lemma that shows that nonzero vectors in  $\mathbb{R}^N$  cannot be mapped to  $\mathcal{Z}$ , the nullspace of diag $(A_e)$ , by the connectivity matrix.

LEMMA 3.3. If  $\mathbf{x} \in \mathbb{R}^N$  is a nonzero vector then  $L\mathbf{x} \notin \mathcal{Z}$ , where L, diag $(A_e)$ , diag $(P_e)$  and  $\mathcal{Z}$  are defined by (2.3), (2.2), (3.2) and (3.6), respectively.

*Proof.* Both  $\mathcal{A}$  and  $\mathcal{P}_{BD}$  are positive definite, which implies that for any  $x \neq 0$ ,

$$\mathcal{A}\boldsymbol{x} = L^T \operatorname{diag}(A_e)(L\boldsymbol{x}) \neq 0 \text{ and } \mathcal{P}_{BD}\boldsymbol{x} = L^T \operatorname{diag}(P_e)(L\boldsymbol{x}) \neq 0.$$

We know from Lemma 3.1 that  $\operatorname{null}(\operatorname{diag}(P_e)) \subset \operatorname{null}(\operatorname{diag}(A_e)) = \mathcal{Z}$ , so  $Lx \notin \mathcal{Z}$ .  $\square$ Both  $\mathcal{A}$  and  $\mathcal{P}_{BD}$  are positive definite and so  $\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A})$  has the variational characterisation [23, Chapters 1 and 15]

$$\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A}) = \min_{\boldsymbol{x} \neq 0} \frac{\boldsymbol{x}^T \mathcal{A} \boldsymbol{x}}{\boldsymbol{x}^T \mathcal{P}_{BD} \boldsymbol{x}} = \min_{\substack{\boldsymbol{y} = L\boldsymbol{x}, \\ \boldsymbol{x} \neq 0}} \frac{\boldsymbol{y}^T \operatorname{diag}(A_e) \boldsymbol{y}}{\boldsymbol{y}^T \operatorname{diag}(P_e) \boldsymbol{y}}.$$

Let  $y = y_R + y_M + y_N$ , where  $y_R \in \mathcal{R}$ ,  $y_N \in \mathcal{N}$ , and  $y_M \in \mathcal{M}$  with  $\mathcal{R}$ ,  $\mathcal{N}$  and  $\mathcal{M}$  defined in (3.6). Lemma 3.3 shows that  $y_R \neq 0$  and so

$$\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A}) = \min_{\substack{\boldsymbol{y} = \boldsymbol{y}_R + \boldsymbol{y}_M \\ \boldsymbol{y}_R \neq \boldsymbol{0}}} \frac{\boldsymbol{y}_R^T \operatorname{diag}(A_e) \boldsymbol{y}_R}{(\boldsymbol{y}_R + \boldsymbol{y}_M)^T \operatorname{diag}(P_e)(\boldsymbol{y}_R + \boldsymbol{y}_M)}.$$
 (3.7)

This appears problematic because, without any restriction on the size of  $\boldsymbol{y}_M$ , the smallest eigenvalue  $\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A})$  could asymptotically tend to zero. To prevent this, we must somehow bound the size of the denominator of (3.7). This is achieved by the next result, provided that  $\boldsymbol{y}_M^T \operatorname{diag}(P_e) \boldsymbol{y}_M < \boldsymbol{y}_R^T \operatorname{diag}(P_e) \boldsymbol{y}_R$ , a condition that we verified numerically for the case of regular grids.

LEMMA 3.4. Let y = Lx,  $x \in \mathbb{R}^N$ ,  $x \neq 0$  be decomposed as

$$\boldsymbol{y} = \boldsymbol{y}_R + \boldsymbol{y}_M + \boldsymbol{y}_N, \tag{3.8}$$

where  $\mathbf{y}_R \in \mathcal{R}$ ,  $\mathbf{y}_N \in \mathcal{N}$  and  $\mathbf{y}_M \in \mathcal{M}$  with  $\mathcal{R}$ ,  $\mathcal{N}$  and  $\mathcal{M}$  defined in (3.6). Additionally, assume that

$$\boldsymbol{y}_{M}^{T}\operatorname{diag}(P_{e})\boldsymbol{y}_{M} \leq \boldsymbol{y}_{R}^{T}\operatorname{diag}(P_{e})\boldsymbol{y}_{R}.$$
 (3.9)

Then

$$(\boldsymbol{y}_R + \boldsymbol{y}_M)^T \operatorname{diag}(P_e)(\boldsymbol{y}_R + \boldsymbol{y}_M) \le 4\boldsymbol{y}_R^T \operatorname{diag}(P_e)\boldsymbol{y}_R.$$
 (3.10)

*Proof.* From Lemma 3.3 we know that  $\mathbf{y}_R \neq \mathbf{0}$ . Since  $\operatorname{diag}(P_e)$  is symmetric positive semidefinite it has a semidefinite square root and there are vectors  $\mathbf{a} \in \mathbb{R}^{16M}$  and  $\mathbf{b} \in \mathbb{R}^{16M}$  for which  $(\mathbf{y}_R + \mathbf{y}_M)^T \operatorname{diag}(P_e)(\mathbf{y}_R + \mathbf{y}_M) = (\mathbf{a} + \mathbf{b})^T (\mathbf{a} + \mathbf{b})$ .

Now, for any vectors  $\boldsymbol{a}$  and  $\boldsymbol{b}$  of the same dimension

$$0 \le \|\boldsymbol{a} - \boldsymbol{b}\|_{2}^{2} = (\boldsymbol{a} - \boldsymbol{b})^{T} (\boldsymbol{a} - \boldsymbol{b}) = 2(\boldsymbol{a}^{T} \boldsymbol{a} + \boldsymbol{b}^{T} \boldsymbol{b}) - (\boldsymbol{a} + \boldsymbol{b})^{T} (\boldsymbol{a} + \boldsymbol{b})$$

or  $(\boldsymbol{a} + \boldsymbol{b})^T (\boldsymbol{a} + \boldsymbol{b}) \le 2(\boldsymbol{a}^T \boldsymbol{a} + \boldsymbol{b}^T \boldsymbol{b})$ . Thus,

$$(\boldsymbol{y}_R + \boldsymbol{y}_M)^T \operatorname{diag}(P_e)(\boldsymbol{y}_R + \boldsymbol{y}_M) \le 2(\boldsymbol{y}_R^T \operatorname{diag}(P_e)\boldsymbol{y}_R + \boldsymbol{y}_M^T \operatorname{diag}(P_e)\boldsymbol{y}_M).$$
 (3.11)

Combining (3.11) with (3.9) gives (3.10).  $\square$ 

We have been unable to prove that (3.9) holds for all meshes for the Dirichlet biharmonic problem, since it does not appear straightforward to remove the influence of the connectivity matrix L. However, there is strong numerical evidence to suggest that the assertion holds. In particular, let  $P_R$  and  $P_M$  be orthogonal projectors onto  $\mathcal{R}$  and  $\mathcal{M}$ , respectively. Then for any vector  $\mathbf{y} = L\mathbf{x}$ ,  $\mathbf{x} \neq \mathbf{0}$ ,

$$\frac{\boldsymbol{y}_R^T\operatorname{diag}(P_e)\boldsymbol{y}_R}{\boldsymbol{y}_M^T\operatorname{diag}(P_e)\boldsymbol{y}_M} = \frac{\boldsymbol{x}^TL^TP_R^T\operatorname{diag}(P_e)P_RL\boldsymbol{x}}{\boldsymbol{x}^TL^TP_M^T\operatorname{diag}(P_e)P_ML\boldsymbol{x}} \geq \delta_{\min}$$

where

$$\delta_{\min} = \lambda_{\min}((L^T P_M^T \operatorname{diag}(P_e) P_M L)^{-1} (L^T P_R^T \operatorname{diag}(P_e) P_R L)). \tag{3.12}$$

The value of  $\delta_{\min}$  is tabulated for different uniform meshes in Table 3.1 and we see that  $\delta_{\min} > 1$  and, asymptotically,  $\delta_{\min}$  appears to tend to 1.05. Thus, (3.9) is satisfied for uniformly refined meshes of square elements.

With these results in hand, we now proceed to bound the smallest eigenvalue of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$ . Under the assumption (3.9), we combine the decomposition (3.8) with Lemma 3.4 to give that, for any y = Lx,  $x \neq 0$ ,

$$\frac{\boldsymbol{y}^T \operatorname{diag}(A_e) \boldsymbol{y}}{\boldsymbol{y}^T \operatorname{diag}(P_e) \boldsymbol{y}} = \frac{\boldsymbol{y}_R^T \operatorname{diag}(A_e) \boldsymbol{y}_R}{(\boldsymbol{y}_R + \boldsymbol{y}_M)^T \operatorname{diag}(P_e) (\boldsymbol{y}_R + \boldsymbol{y}_M)} \ge \frac{1}{4} \frac{\boldsymbol{y}_R^T \operatorname{diag}(A_e) \boldsymbol{y}_R}{\boldsymbol{y}_R^T \operatorname{diag}(P_e) \boldsymbol{y}_R}, \quad (3.13)$$

where by Lemma 3.3,  $\boldsymbol{y}_R \neq \boldsymbol{0}$ . It follows from (3.7) that

$$\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A}) \geq rac{1}{4} \min_{\substack{oldsymbol{y}_R \in \mathcal{R} \ oldsymbol{y}_R 
eq oldsymbol{y}_R \ oldsymbol{y}_R}} rac{oldsymbol{y}_R^T \operatorname{diag}(A_e)oldsymbol{y}_R}{oldsymbol{y}_R^T \operatorname{diag}(P_e)oldsymbol{y}_R}.$$

Since diag $(P_e)$ , diag $(A_e)$  and the projector  $P_R$  onto  $\mathcal{R}$  are block diagonal, the above minimisation over all nonzero  $\boldsymbol{y}_R$  can be carried out using individual element matrices. We computed the minimum for our element matrices and found that for all nonzero  $\boldsymbol{y}_R \in \mathcal{R}$ 

$$\frac{\boldsymbol{y}_R^T \operatorname{diag}(A_e) \boldsymbol{y}_R}{\boldsymbol{y}_R^T \operatorname{diag}(P_e) \boldsymbol{y}_R} > 0.046.$$

Thus,

$$\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A}) \geq \frac{1}{4} \min_{\boldsymbol{y}_R \neq \boldsymbol{0}} \frac{\boldsymbol{y}_R^T \operatorname{diag}(A_e) \boldsymbol{y}_R}{\boldsymbol{y}_R^T \operatorname{diag}(P_e) \boldsymbol{y}_R} > 0.0115.$$

Combining (3.9) with Theorem 3.2 gives the following bounds on the eigenvalues of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$ .

COROLLARY 3.5. Let  $\mathcal{A}$  and  $\mathcal{P}_{BD}$  be as in (1.4) and (3.1) and assume that (3.9) holds. Then the eigenvalues  $\lambda$  of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  satisfy 0.0115  $< \lambda \le 2$  and  $\kappa_2(\mathcal{P}_{BD}^{-1}\mathcal{A}) < 174$  independently of the mesh spacing parameter h.

Comparison with Table 3.1 shows that the bounds in Corollary 3.5 are pessimistic. However, combined with the high multiplicity of the unit eigenvalue, they show that we can expect fast convergence whenever (3.9) is satisfied.

Table 3.1

Smallest  $(\lambda_{\min})$  and largest  $(\lambda_{\max})$  eigenvalues of the preconditioned operator  $\mathcal{P}_{BD}^{-1}\mathcal{A}$ ,  $r = \operatorname{rank}(B)$  from Theorem 3.2 and  $\delta_{\min}$  from (3.12) for a sequence of uniformly refined grids.

Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$
$\lambda_{\min}$	0.72	0.64	0.61	0.60	0.60
$\lambda_{ m max}$	1.28	1.36	1.39	1.40	1.40
rank(B)	9	49	225	961	3969
$\delta_{\min}$	1.17	1.06	1.05	1.05	1.05

4. A practical preconditioner. Although the preconditioner  $\mathcal{P}_{BD}$  in (3.1) has favourable spectral properties, it is prohibitively expensive to apply for large problems, since it requires the solution of linear systems involving the  $3n \times 3n$  matrix A. We will now investigate the performance of the block bordered diagonal preconditioner

$$\mathcal{P}_{BBD} = \begin{bmatrix} A_{11} & A_{12} & A_{13} \\ A_{12}^T & A_{22} \\ A_{13}^T & A_{33} \\ & & A_{44} \end{bmatrix}$$
(4.1)

that is formed by omitting the blocks  $A_{23}$  and  $A_{23}^T$  from  $\mathcal{P}_{BD}$ . Unlike  $\mathcal{P}_{BD}$ , the symmetric positive definiteness of  $\mathcal{A}$  is not enough to guarantee that  $\mathcal{P}_{BBD}$  is positive definite, but computations show that it is positive definite for all grids tested. Compared to the even simpler block Jacobi preconditioner

$$\mathcal{P}_{J} = \begin{bmatrix} A_{11} & & & & \\ & A_{22} & & & \\ & & A_{33} & & \\ & & & A_{44} \end{bmatrix}, \tag{4.2}$$

 $\mathcal{P}_{BBD}$  retains the coupling between u and both first derivative  $(\frac{\partial u}{\partial s_1})$  and  $\frac{\partial u}{\partial s_2}$  DOFs. We will see in Section 5 that retaining this coupling is essential to obtaining low and asymptotically constant iteration counts as the mesh is refined.

The action of  $\mathcal{P}_{BBD}^{-1}$  on a vector can be computed by means of the unsymmetric UL decomposition

$$\mathcal{P}_{BBD} = UL = \begin{bmatrix} I & A_{12}A_{22}^{-1} & A_{13}A_{33}^{-1} \\ & I & \\ & & I \end{bmatrix} \begin{bmatrix} S_{11} \\ A_{12}^T & A_{22} \\ A_{13}^T & & A_{33} \\ & & & A_{44} \end{bmatrix}, \quad (4.3)$$

where

$$S_{11} = A_{11} - A_{12}A_{22}^{-1}A_{12}^{T} - A_{13}A_{33}^{-1}A_{13}^{T}. (4.4)$$

Note that the solve Uw = v can be performed in a block parallel manner.

The remainder of this section is devoted to understanding the spectral properties of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  and deriving an approximation that can be implemented in a cost-optimal manner.

**4.1. Eigenvalue analysis.** The block structure of  $\mathcal{P}_{BBD}$  and, in particular, the indefiniteness of the element matrices used in its assembly prevent us from applying

the previously introduced analysis to bound the spectrum of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$ . Instead, we consider the eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD}$  and then apply the bounds

$$\lambda_{\min}(\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD})\lambda_{\min}(\mathcal{P}_{BD}^{-1}\mathcal{A}) \leq \lambda(\mathcal{P}_{BBD}^{-1}\mathcal{A}) \leq \lambda_{\max}(\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD})\lambda_{\max}(\mathcal{P}_{BD}^{-1}\mathcal{A}),$$
(4.5)

which follow from the Courant-Fischer theorem [17, Theorem 4.2.11], in conjunction with the bounds in Corollary 3.5. The eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD}$  are given in the following lemma.

LEMMA 4.1. Let rank $(A_{23}) = s$ . Then 1 is an eigenvalue of  $\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD}$  with multiplicity N-2s while the remaining 2s eigenvalues  $\eta$  satisfy

$$(G - FA_{11}^{-1}F^{T})\mathbf{v} = \eta(\widetilde{G} - FA_{11}^{-1}F^{T})\mathbf{v}, \tag{4.6}$$

where  $\mathbf{v} \neq \mathbf{0}$ ,

$$F^{T} = \begin{bmatrix} A_{12} & A_{13} \end{bmatrix}, \quad G = \begin{bmatrix} A_{22} & A_{23} \\ A_{23}^{T} & A_{33} \end{bmatrix} \text{ and } \widetilde{G} = \begin{bmatrix} A_{22} \\ A_{33} \end{bmatrix}.$$
 (4.7)

*Proof.* In the notation of (3.3)

$$\mathcal{P}_{BBD} = \begin{bmatrix} \widetilde{A} & \\ & A_{44} \end{bmatrix}, \quad \widetilde{A} = \begin{bmatrix} A_{11} & A_{12} & A_{13} \\ A_{12}^T & A_{22} & \\ A_{13}^T & & A_{33} \end{bmatrix}$$

and

$$\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD} = \begin{bmatrix} \widetilde{A}^{-1}A & \\ & I_n \end{bmatrix},$$

where  $I_n$  is the identity matrix of dimension n. This shows that 1 is an eigenvalue of  $\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD}$  with multiplicity at least n.

To obtain the remaining 3n eigenvalues let us further partition  $\widetilde{A}$  and A as

$$\widetilde{A} = \begin{bmatrix} A_{11} & F^T \\ F & \widetilde{G} \end{bmatrix}, \qquad A = \begin{bmatrix} A_{11} & F^T \\ F & G \end{bmatrix},$$

where F and G are as in (4.7). Then, the result is obtained by a straightforward extension of Theorem 3.1 of Dollar et al. [10] to the case of rank-deficient F, which we sketch out for completeness.

The eigenvalues  $\eta$  of  $\tilde{A}^{-1}A$  satisfy

$$A_{11}\boldsymbol{u} + F^T\boldsymbol{v} = \eta A_{11}\boldsymbol{u} + \eta F^T\boldsymbol{v}, \tag{4.8}$$

$$F\boldsymbol{u} + G\boldsymbol{v} = \eta F\boldsymbol{u} + \eta \widetilde{G}\boldsymbol{v},\tag{4.9}$$

where  $\boldsymbol{u} \in \mathbb{R}^{2n}$  and  $\boldsymbol{v} \in \mathbb{R}^n$  are not simultaneously zero. From (4.8) we see that either  $\eta = 1$  or  $A_{11}\boldsymbol{u} + F^T\boldsymbol{v} = \boldsymbol{0}$ . If  $\eta = 1$  then, letting  $\boldsymbol{v} = [\boldsymbol{v}_1^T \ \boldsymbol{v}_2^T]^T$  with  $\boldsymbol{v}_1, \boldsymbol{v}_2 \in \mathbb{R}^n$ , we find that there are n-s linearly independent vectors  $\boldsymbol{v}_1 \in \text{null}(A_{23})$  for which (4.8) and (4.9) are satisfied with  $\boldsymbol{v}_2 = \boldsymbol{u} = \boldsymbol{0}$ . Similarly, there are n-s linearly independent vectors  $\boldsymbol{v}_2 \in \text{null}(A_{23}^T)$  for which (4.8) and (4.9) are satisfied with  $\boldsymbol{v}_1 = \boldsymbol{u} = \boldsymbol{0}$ . Otherwise,  $\boldsymbol{v} = \boldsymbol{0}$  and we can find  $\boldsymbol{n}$  linearly independent vectors  $\boldsymbol{u} \neq \boldsymbol{0}$ . Combining these results shows that  $\eta = 1$  with multiplicity 3n-2s. If  $\eta \neq 1$  then  $\boldsymbol{u} = -A_{11}^{-1}F^T\boldsymbol{v}$  and substituting into (4.9) gives (4.6).  $\square$ 

Table 4.

Computed smallest  $(\lambda_{\min})$  and largest  $(\lambda_{\max})$  eigenvalues of the preconditioned operator  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  as well as  $s=\operatorname{rank}(A_{23})$  for a sequence of uniformly refined grids.

Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$
$\overline{N}$	36	196	900	3844	15876
$\lambda_{\min}$	0.72	0.62	0.58	0.56	0.55
$\lambda_{ m max}$	1.27	1.38	1.39	1.41	1.41
$rank(A_{23})$	8	48	224	960	3968

Similarly to Theorem 3.2 we see that  $\lambda = 1$  is an eigenvalue of  $\mathcal{P}_{BBD}^{-1}\mathcal{P}_{BD}$  with high multiplicity. However, we have been unable to bound the remaining 2s eigenvalues. In spite of this, combining Lemma 4.1 with Corollary 3.5 and (4.5) shows that most of the eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  lie in a bounded interval.

COROLLARY 4.2. Whenever (3.9) is satisfied at least N-2s eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  lie in (0.0115,2) and the remaining eigenvalues lie in (0.0115 $\eta$ ,2 $\eta$ ), where  $\eta$  is an eigenvalue of the generalised eigenvalue problem (4.6).

The extreme eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  are given in Table 4.1 as a function of the problem size N. From this we see that these eigenvalues do not differ greatly from the extreme eigenvalues of  $\mathcal{P}_{BD}^{-1}\mathcal{A}$  (see Table 3.1), and in practice little is lost in terms of the asymptotic convergence speed by using a more practical preconditioner. The numerical evidence in Table 4.1 suggests that the extreme eigenvalues of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  appear to be bounded under mesh refinement, although we have been unable to prove this analytically.

**4.2. Further simplifications.** Although the block decomposition (4.3) allows the efficient application of  $\mathcal{P}_{BBD}$ , to achieve a preconditioner with optimal cost we require optimal solvers for linear systems involving the principal diagonal blocks  $S_{11}$ ,  $A_{22}$ ,  $A_{33}$  and  $A_{44}$ .

First, we consider spectrally equivalent approximations of  $A_{22}$ ,  $A_{33}$  and  $A_{44}$ . LEMMA 4.3. Let

$$L_{22} = \text{lump}(A_{22}), \quad L_{33} = \text{lump}(A_{33}),$$
 (4.10)

where  $lump(H) = \{h_{ij}\}$  with

$$h_{ij} = \begin{cases} \sum_{k=1}^{n} h_{ik}, & i = j, \\ 0, & i \neq j. \end{cases}$$

Then the eigenvalues of  $L_{22}^{-1}A_{22}$  and  $L_{33}^{-1}A_{33}$  are contained in (1/3,1) while the eigenvalues of diag $(A_{44})^{-1}A_{44}$  are contained in (0.43,1.24).

Proof. The matrices  $A_{22}$ ,  $A_{33}$  and  $A_{44}$  are assembled from  $4 \times 4$  submatrices  $A_e^{(22)}$ ,  $A_e^{(33)}$  and  $A_e^{(44)}$  of the element matrix  $A_e$ . Additionally, the approximations  $L_{22}$  and  $L_{33}$  are assembled from lumped versions of  $A_e^{(22)}$  and  $A_e^{(33)}$ , while diag $(A_{44})$  is assembled from the diagonal of  $A_e^{(44)}$ . All six of these element matrices are SPD. As a result, we can use the approach of Wathen [27] to prove the result. (Recall that a similar result was used in Section 3, where we had to deal with singular element matrices.)  $\square$ 

Remark 1. This result is not surprising since the spectrum of  $A_{22} = A_{33}$  resembles that of a scaled mass matrix, and for such matrices lumping often gives spectrally equivalent operators. In particular, for our problem  $\lambda(A_{22}) = \lambda(A_{33}) \sim O(h^{-4})\lambda(M)$ ,

where M is a mass matrix; for uniform grids this can be verified using Fourier analysis, similarly to the approach in [11, Section 1.6]. Additionally, on a uniform mesh all entries of  $A_{22}$  and  $A_{33}$  are nonnegative. However, this nonnegativity is lost for some stretched grids, the result of which is a deterioration of the convergence rate (see Section 5.1).

Using (4.3) and Lemma 4.3 we approximate  $\mathcal{P}_{BBD}$  by

$$\widetilde{\mathcal{P}}_{BBD} = \begin{bmatrix} I & A_{12}L_{22}^{-1} & A_{13}L_{33}^{-1} \\ I & & & \\ & I & & \\ & & & I \end{bmatrix} \begin{bmatrix} \overline{S}_{11} & & & \\ A_{12}^T & L_{22} & & & \\ A_{13}^T & & L_{33} & & \\ & & & & \operatorname{diag}(A_{44}) \end{bmatrix}, \quad (4.11)$$

where

$$\overline{S}_{11} = A_{11} - A_{12}L_{22}^{-1}A_{12}^{T} - A_{13}L_{33}^{-1}A_{13}^{T}. \tag{4.12}$$

The block  $\overline{S}_{11}$  in (4.12) is a sparse approximation of the Schur complement  $S_{11}$  from (4.4), owing to the diagonal approximations (4.10), and can be assembled cheaply.

To apply the preconditioner  $\widetilde{\mathcal{P}}_{BBD}$  within the preconditioned CG algorithm we must solve systems with  $\overline{S}_{11}$ ,  $L_{22}$ ,  $L_{33}$  and diag( $A_{44}$ ). The last three matrices are diagonal, and hence trivial to invert. For systems with  $\overline{S}_{11}$  we consider two approaches: an LU factorisation, which yields an exact solution but is not computationally optimal, or a small, fixed number of V-cycles of classical algebraic multigrid (AMG) [20, p. 73] with various smoothers. The latter has optimal cost but leads to an inexact solution. We found that the application of two V(2,2) cycles of AMG with Gauss-Seidel smoothing and Ruge-Stüben coarsening [20] as implemented in the HSL routine MI\_20 [2, 18] was most effective in this context for small problems, while for larger problems we used ILU(0) smoothing as implemented in BoomerAMG from the Hypre library. More details on the AMG method used are given in Section 5.

Using these approximations for the Schur complement subsystem, we obtain the preconditioners  $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$  and  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  in which the Schur complement subsidiary system is solved using an LU factorisation and AMG, respectively. In Table 4.2 we present the spectral properties of the preconditioned operators  $\left(\widetilde{\mathcal{P}}_{BBD}^{[LU]}\right)^{-1}\mathcal{A}$  and  $\left(\widetilde{\mathcal{P}}_{BBD}^{[AMG]}\right)^{-1}\mathcal{A}$ . These results suggest that the spectrum of  $\left(\widetilde{\mathcal{P}}_{BBD}^{[LU]}\right)^{-1}\mathcal{A}$  is bounded under mesh refinement, as we might expect from the spectral equivalence bounds in Lemma 4.3, although the eigenvalues are not as tightly clustered as those of  $\mathcal{P}_{BBD}^{-1}\mathcal{A}$  in Table 3.1. However, the smallest eigenvalue of  $\left(\widetilde{\mathcal{P}}_{BBD}^{[AMG]}\right)^{-1}\mathcal{A}$  decreases with mesh refinement; that is, the AMG approximation with a Gauss-Seidel smoother is not spectrally equivalent to  $\overline{S}_{11}$ .

5. Numerical experiments. In this section we examine the effectiveness of the preconditioners  $\mathcal{P}_{BD}$ ,  $\mathcal{P}_{BBD}$  and  $\widetilde{\mathcal{P}}_{BBD}$  at reducing the number of conjugate gradient iterations and the computational time. Additionally, we investigate the robustness of their performance with respect to deformations of the domain and stretching of the finite elements. Throughout, we choose the homogeneous Dirichlet boundary conditions  $g_1 = g_2 = 0$  in (1.2). Unless otherwise specified we consider a unit square domain  $\Omega = [0, 1]^2$  discretised by a uniform grid of square elements. Although we note that for finite element problems the stopping criterion for CG should be tied to

Table 4.2 Smallest ( $\lambda_{min}$ ) and largest ( $\lambda_{max}$ ) eigenvalues of  $\left(\widetilde{\mathcal{P}}_{BBD}^{[LU]}\right)^{-1}\mathcal{A}$  and  $\left(\widetilde{\mathcal{P}}_{BBD}^{[AMG]}\right)^{-1}\mathcal{A}$  for a sequence of uniformly refined grids. Note that we were unable to obtain the eigenvalues of the largest  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  preconditioned matrix because of memory constraints.

	Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$
$(\widetilde{\mathcal{D}}^{[LU]})^{-1}$	$\lambda_{\min}$	0.40	0.33	0.30	0.29	0.28
$(P_{BBD})$ A	$\lambda_{ m max}$	1.25	1.30	1.31	1.32	1.32
$\overline{\left(\widetilde{\mathcal{P}}_{BBD}^{[AMG]} ight)^{-1}\mathcal{A}}$	$\lambda_{\min}$	0.40	0.31	0.21	0.13	
$\left(\mathcal{P}_{BBD}^{[BBD]}\right)$ $\mathcal{A}$	$\lambda_{ m max}$	1.25	1.30	1.31	1.32	_

the discretisation error, to demonstrate mesh independence we terminate the preconditioned CG method when the residual decreases in norm by six orders of magnitude, that is,  $\|\boldsymbol{r}^{(k)}\|_2 \leq 10^{-6}\|\boldsymbol{r}^{(0)}\|_2$ .

We first compare preconditioned CG iterations for  $\mathcal{P}_{BD}$ ,  $\mathcal{P}_{BD}$   $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$  and  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$ for smaller problems using Matlab. For comparison, we also present iteration counts for the block Jacobi preconditioner (4.2) and a black-box AMG preconditioner applied to the entire matrix  $\mathcal{A}$ . The AMG method we use is the HSL code MI20 [2, 18], with two-pass Ruge-Stüben coarsening and point Gauss-Seidel smoothing. We select the default options except that we change the coarsening criterion c\_fail from 1 to 2 and use two V(2,2) cycles. We note that different AMG methods may give different results. However, our aim is to develop an effective preconditioner that is easy to implement, and so we choose an off-the-shelf code that generally works well for finite element problems [2]. We stress that preconditioners  $\mathcal{P}_{BBD}$  in (4.3) and  $\widetilde{\mathcal{P}}_{BBD}$  in (4.11) are parallelisable, like the block Jacobi preconditioner  $\mathcal{P}_{J}$  in (4.2), as discussed in Section 4. Since the problems considered here are of relatively small dimension, in addition to measuring the norm of the residual we compute the relative error  $\|\boldsymbol{x} - \boldsymbol{x}^{(k)}\|_{\mathcal{A}} / \|\boldsymbol{x}\|_{\mathcal{A}}$  in the energy norm at termination. In all cases this is smaller than  $1.6 \times 10^{-6}$ . Computations were performed with different right-hand-sides **b**: we used the right-hand-side from the finite element discretisation of (1.1) for f=1 and a random right-hand side b. Both choices give similar results, which shows that the convergence results do not depend on the regularity of the forcing term. Consequently, only results for f = 1 are presented.

The results are given in Table 5.1, from which we see that without preconditioning the number of CG iterations increases rapidly, and appears to grow as  $O(h^{-2})$ . The application of the AMG and block Jacobi preconditioners reduces the number of iterations somewhat, but convergence is still mesh dependent. This is not surprising since the condition numbers of these preconditioned matrices increase as the mesh is refined (see Table 5.2). However, the number of iterations required for  $\mathcal{P}_{BD}$ ,  $\mathcal{P}_{BBD}$  and  $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$  seem to be bounded independently of the mesh. This is in line with the spectral analysis in previous sections and the computed eigenvalues in Tables 3.1, 4.1 and 4.2.

To explore the asymptotic behaviour of  $\widetilde{\mathcal{P}}_{BBD}$  for larger problems, Figures 5.1 and 5.2 show the number of iterations required for convergence of preconditioned CG, and the solution times, for the three block bordered diagonal preconditioners ( $\mathcal{P}_{BBD}$ , and  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$ ). These results were obtained using a C++ implementation with HYPRE's BOOMERAMG [16] for the AMG solver with V(2,2) cycles and SUPERLU [9] for the direct solver. Note that in the case of  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  we used an ILU(0) smoother in these experiments instead of a Gauss-Seidel smoother (an option that is available in

Table 5.1

CG iteration counts for the unpreconditioned system, and preconditioned CG iterations counts for several different preconditioners: AMG applied to the whole matrix  $\mathcal{A}$  in (1.4),  $\mathcal{P}_J$ ,  $\mathcal{P}_{BD}$ ,  $\mathcal{P}_{BBD}$  and the two inexact versions  $\mathcal{P}_{BBD}^{[LU]}$ , and  $\mathcal{P}_{BBD}^{[AMG]}$ .

Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$	$128 \times 128$
N	36	196	900	3844	15876	64516
Unpreconditioned	22	52	172	640	2459	9742
AMG	3	9	27	101	381	1378
$\mathcal{P}_J$	13	35	80	168	339	678
$\mathcal{P}_{BD}$	7	9	9	10	10	10
$\mathcal{P}_{BBD}$	7	9	10	11	11	11
$\widetilde{\mathcal{P}}_{BBD}^{[LU]}$	11	14	14	15	16	16
$\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$	11	14	17	22	29	40

Table 5.2

Smallest  $(\lambda_{\min})$  and largest  $(\lambda_{\max})$  eigenvalues of the AMG and Jacobi preconditioned matrices. Note that we were unable to obtain the eigenvalues of the largest AMG preconditioned matrix because of memory constraints.

	Elements	$4 \times 4$	$8 \times 8$	$16 \times 16$	$32 \times 32$	$64 \times 64$
$\mathcal{D}^{-1}$ $\Lambda$	$\lambda_{\min}$	0.92	0.17	0.01	0.0008	
$\mathcal{P}_{AMG}^{-1}\mathcal{A}$	$\lambda_{ m max}$	1.00	1.00	1.00	1.00	_
$\mathcal{P}_{I}^{-1}\mathcal{A}$	$\lambda_{\min}$	0.18	0.04	0.009	0.002	0.0005
$P_J^- \mathcal{A}$	$\lambda_{ m max}$	1.80	2.02	2.07	2.10	2.10

BoomerAMG but not in HSL\_MI20). Computations were performed on a 3.6 GHz Intel Xeon processor.

Figure 5.1 shows that the iteration count increases when  $\mathcal{P}_{BBD}$  is replaced by  $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$ , and the inexact version  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$ . For all three preconditioners the number of iterations appears to be bounded independently of the mesh parameter h. Note that it was not possible to obtain results for  $\mathcal{P}_{BBD}$  for the larger problems as a result of memory limitations.

Figure 5.2 shows the solution times for the different strategies. For comparison, we also include the execution time when SUPERLU is used to solve the full linear system (1.4). We see that all the preconditioned CG solutions are obtained significantly faster than with the direct method SUPERLU and that the approximations in  $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$  and  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  significantly reduce the execution time. However, only preconditioned CG with  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  yields a solution that scales approximately linearly with problem size; this is because the number of iterations appears to be bounded when N increases and the cost of applying the preconditioner itself scales linearly with the problem size. The other variants do not possess this property because they involve computationally sub-optimal LU factorisations.

**5.1. Robustness of the preconditioner.** The block bordered diagonal preconditioner  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  with ILU(0) smoothing was shown to be optimal in terms of wall clock time for a simple test problem. A key step in the development of  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  was the approximation of the blocks  $A_{22}$  and  $A_{33}$  by their lumped counterparts  $L_{22}$  and  $L_{33}$ , and the application of AMG to the Schur complement  $\overline{S}_{11}$  in (4.11). However, AMG is known to be less effective on problems with strongly stretched or distorted elements [8], and the lumped matrices can be poor approximations in this case (cf. Remark 1). Therefore we will now evaluate the robustness of our preconditioner for problems with stretched grids and non-square domains and elements. Stretched grids

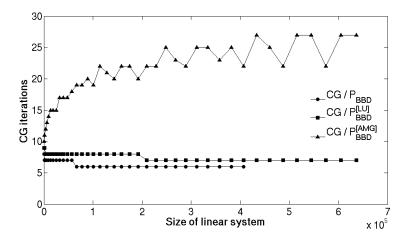


Fig. 5.1. Number of preconditioned CG iterations for approximations of the  $\mathcal{P}_{BBD}$  preconditioner.

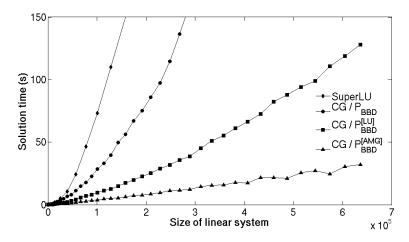


Fig. 5.2. Solution times (in seconds) for CG, preconditioned with approximations of the  $\mathcal{P}_{BBD}$  preconditioner. The execution time of the direct method SuperLU applied to the system with coefficient matrix (1.4) is presented for comparison.

are needed, for example, for accurate computations of biharmonic eigenfunctions near the corners of the domain (see [5]). Figures 5.3 and 5.4 illustrate the two tests considered:

- 1. **Stretched Elements.** The domain is stretched in the  $x_1$  direction and the deformation is described by the aspect ratio a.
- 2. **Distorted Elements.** The top right corner is stretched upwards in the  $x_2$  direction; the ratio of the heights of the vertical boundaries is parameterised by b.

In both tests we used the same number of elements in each coordinate direction.

Figures 5.5 and 5.6 show the iteration counts for these two tests. We find that the iteration count increases with an increasing degree of deformation, but it does not appear to grow significantly as the mesh is refined. Since the loss of preconditioning

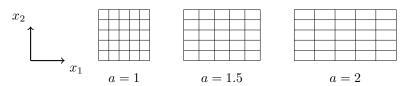


Fig. 5.3. Robustness test 1: Stretched elements. The domain is stretched in the  $x_1$  direction and the deformation is described by the aspect ratio a.

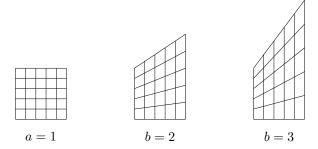


Fig. 5.4. Robustness test 2: Distorted domain and elements. The top right corner is stretched upwards in the  $x_2$  direction; the ratio of the heights of the vertical boundaries is parameterised by b.

performance is partly due to the use of AMG to solve the Schur complement subsidiary linear system, replacing AMG with an LU factorisation (i.e. using  $\widetilde{\mathcal{P}}_{BBD}^{[LU]}$ ) allows more strongly deformed problems to be solved. However, eventually this preconditioner also starts to lose its effectiveness. This is because with small grid stretching or grid/domain deformations the inefficiency comes from the algebraic multigrid method, but as the stretch and deformation rates increase, problems with the approximations of  $A_{22}$ ,  $A_{33}$  and  $A_{44}$  also contribute to increased iteration counts.

6. Conclusions. We have presented effective preconditioners for the  $C^1$  finite element discretisation of the Dirichlet biharmonic problem using Hermitian bicubic elements. The preconditioners are easy to set up as they only involve operations on blocks that are readily extracted from the full system; these blocks can also be computed from element matrices. On uniform meshes both the block diagonal and block bordered diagonal preconditioners appear to give mesh independent convergence. Moreover, we analysed the spectrum of block diagonal and block bordered diagonal preconditioners and showed that, under a certain condition, the block diagonal preconditioner  $\mathcal{P}_{BD}$  gives mesh independent convergence; the required condition holds for all the uniform meshes tested here. Our analysis of the block diagonal preconditioner uses the approach of Wathen [27], [28], [29], which assumes that the coefficient matrix and preconditioner are symmetric positive definite and are assembled from individual element matrices. As such, Wathen's appealing technique is applicable to other finite element discretisations, differential operators and preconditioners.

To obtain a cost-optimal implementation, we further simplified the block bordered diagonal preconditioner  $\mathcal{P}_{BBD}$  by lumping certain block matrices and using AMG for the approximate solution of a Schur complement subsidiary linear system. This preconditioner also gave iteration counts for uniform grids that seemed to level off as the mesh was refined. We tested this approximate preconditioner on stretched and distorted elements and found that, although stretching and distorting elements

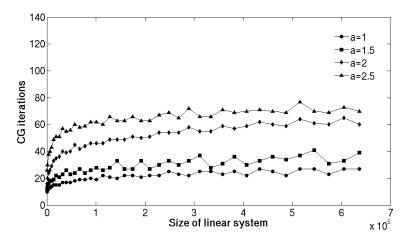


Fig. 5.5. Number of CG iterations for  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  in robustness test 1 (stretched elements).

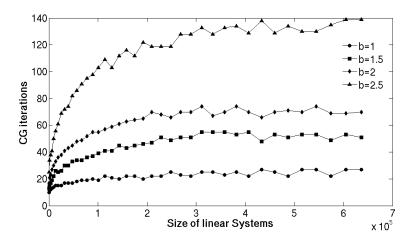


Fig. 5.6. Number of CG iterations for  $\widetilde{\mathcal{P}}_{BBD}^{[AMG]}$  in robustness test 2 (distorted elements).

increased the number of iterations, for moderately deformed elements this increase was mild, and the number of iterations did not grow significantly as the mesh was refined.

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